SE 4C03 Winter 2002

11. Routing Protocols

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Routing Management Goals

- Packets can get from any source to any destination
- Packets take the shortest route
- Routing information is automatically distributed
- Routing tables are automatically initialized and updated
- Many routers route with partial information about destinations by using default routes
  - Tables are smaller and routing is more efficient
  - Management can be performed locally
  - Routing tables may be inconsistent with each other

Core Router Architecture

- Pool of internet routers are divided into **core routers** and **noncore routers**
  - Motivated by the ARPANET which was a single WAN
  - Works well when the internet has a single backbone
- Provides optimal routes for all possible destinations
  - Core routers do not use default routes
  - Every assigned class network address must be advertised to the core system
  - A noncore router sends all nonlocal traffic to the core router at its site
  - Core routers communicate with themselves to preserve consistency
- Core router architecture is impractical today because:
  - The Internet is no longer built around one backbone
  - It is not possible to have a core router at each site
  - Core router architecture does not scale up very well

Peer Backbone Architecture

- **Peer backbone networks** have several connections with each other
  - Motivated by connection of the NSFNET backbone to the ARPANET backbone
- The system is difficult to implement because:
  - Packets cannot simply be routed according to the network portion of their destination
  - Peers must keep their routes consistent with each other
  - Defaults routes from one peer to another can create routing loops

Vector Distance Routing (1)

- Each router keeps a list of route records having the following fields:
  - The destination class network (i.e., the **vector**)
  - The number of hops to the destination (i.e., the **distance**)
  - The route (either **direct** or the name of a **router**)
- Assumption: Measuring distance as number of hops is a good measure of the time cost of a route
- When a router boots, the list is initialized to just the routes for the class networks that are directly connected to the router
- Each router sends a copy of its list to all other routers that are directly connected to it

Vector Distance Routing (2)

- Suppose that R_1 and R_2 are directly connected routers
  - If R_1 has a route record (N, D, R) for the class network N but R_2 does not, R_2 will add the record $(N, D+1, R_1)$ to its list
  - If R_1 and R_2 have route records (N, D_1, R_1') and (N, D_2, R_2'), respectively, with $D_1 + 1 < D_2$, then R_2 will update its record to $(N, D_1 + 1, R_1)$
  - If R_1 and R_2 have route records (N, D_1, R) and (N, D_2, R_1), respectively, with $D_1 + 1 \neq D_2$, then R_2 will update its record to $(N, D_1 + 1, R_1)$

Vector Distance Routing (2)

- Advantage: Vector-distance algorithms are easy to implement
- Disadvantage: Vector-distance algorithms do not scale up well
  - Each router eventually has a record in its list for each class network
  - Lots of routing information is transmitted
  - Routing information is propagated slowly (which can make the system unstable)
  - Any incorrect routing information will be propagated along with correct information

Gateway-to-Gateway Protocol (GGP)

- Now defunct vector-distance protocol used for sharing routing information among core routers
- GGP messages were encapsulated in IP datagrams
- Kinds of GGP messages:
  - Routing update message
  - Positive acknowledgement to routing update (update acceptable)
  - Negative acknowledgement to routing update (error detected)
  - Echo request
  - Echo reply

Link-State Routing (1) (Shortest Path First (SPF) Routing)

- Each router keeps a graph of the topology of the internet
  - A node represents a router
  - An edge represents a network
- The routers work to keep the graphs up to date
  - Each router periodically checks to see if its neighbors are up or down
  - Each router periodically broadcasts a message that contains the state of each of its links
  - Each router uses the link-state broadcasts it receives to update its internet topology graph

Link-State Routing (2)

- Whenever the graph changes, a router uses the **Dijkstra shortest path algorithm** to compute the shortest route to each destination
- Advantages over vector-distance routing
  - Each router independently computes the shortest routes from the same information
  - It is easy to fix mistakes because link-state information is not modified as it is propagated
  - The size of link-state messages do not grow as the internet grows

Problems with Core Router System

- It is impractical for the group of core routers to include more than a small portion of the internet routers
- Noncore routers need to know core router routes to avoid "extra hops"
- Core routers need to know about "hidden" networks

Autonomous System Architecture

- An internet is composed of several **autonomous systems** each composed of a collection of routers and networks
- Each autonomous system chooses its own internal routing architecture
  - An **interior gateway protocol (IGP)** is used to distribute routing information within an autonomous system
- One or more routers in an autonomous system advertises local routing information to other autonomous systems
  - An **exterior gateway protocol (EGP)** is used to distribute routing information between two autonomous systems

Border Gateway Protocol (BGP)

- BGP is the EGP currently used in the Internet and most other TCP/IP internets
  - Each autonomous system has at least one designated **border gateway** that speaks on behalf of the autonomous system
  - Border gateways exchange reachability information (but not optimal routes)
- Attributes of BGP:
  - Transport is via TCP and thus is reliable
  - Local routing policy can be supported
  - Updates are usually incremental
  - A receiver can authenticate the sender
  - Subnet addressing is supported and so, for example, related destinations can be aggregated

BGP Message Types

1. OPEN: initializes communication
2. UPDATE: withdraws unreachable destinations and advertises new destinations
   - Each advertised destination can include:
     - The next hop to the destination
     - The path of autonomous systems to the destination
     - The source of the advertisement
   - Each destination is expressed as a compressed mask-address pair
3. NOTIFICATION: reports an error
4. KEEPALIVE: tests network connectivity
   - KEEPALIVE messages have minimum size (they contain only a BGP header)

Routing Information Protocol (RIP) (1)

- An IGP protocol for vector-distance routing
  - Distance is measured by the number of hops
- Implemented by the `routed` program designed at the University of California at Berkeley
  - Became popular because it was distributed with BSD UNIX
- RIP participants are either **active** or **passive**
  - Active participants are routers that advertise their routes to others
  - Passive participants are hosts that do not advertise routes but use advertisements to update their routes

Routing Information Protocol (RIP) (2)

- An active participant broadcasts a message every 30 seconds
  - RIP messages are encapsulated in UDP datagrams
  - A RIP server listens at UDP port 520
- Route restrictions
  - An old route is retained until a new one with a strictly lower cost is received
  - Routes timeout after 180 seconds
  - The maximum possible distance is 16, so RIP only works with relatively small autonomous systems
- RIP suffers from the problems that are inherent in vector-distance routing
- RIP does not employ router authentication

The HELLO Protocol

- A defunct IGP protocol for vector-distance routing
- Routes are measured by network delay instead of number of hops

The Open SPF Protocol (OSPF) (1)

- An IGP protocol for link-state (SPF) routing
- Provides support for:
  - **Type of service routing**: Routing is done on the basis of both destination address and type of service (precedence plus low delay, high throughput, or high reliability)
  - **Load balancing**: The same traffic can be distributed over multiple routes
  - **Area organization**: Local networks and routers can be organized into independent **areas**
  - **Authentication**: Routers must authenticate each other
  - **Subnet routes**: Routes may be directed to subnets
  - **Information forwarding**: Routers may forward information received from routers exterior to a site

The Open SPF Protocol (OSPF) (2)

  - **Graph structure**: Each router keeps a graph of the topology of the internet whose structure is more complex than that of other link-state routing protocols
    - A node represents a router or an SPN
    - One router can be directly connected to another router in the graph
- Kinds of OSPF messages:
  - Hello (to test router reachability)
  - Graph description (to initialize a router's topology graph)
  - Link status (to update a link in a router's topology graph)
- OSPF messages are encapsulated in IP datagrams

The gated Program

- Handles multiple routing protocols including both IGPs and EGPs
- Allows a router to communicate with routers both inside and outside its autonomous system
- IGPs supported include: RIP, HELLO, OSPF
- EGPs supported include: BGP