

SE 4C03 Winter 2007

## 03 Internet Addressing

William M. Farmer

Department of Computing and Software  
McMaster University

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# IP Addresses

- There are two Internet naming systems:
  1. The primary system is the **internet address system** which uses binary **IP addresses**.
  2. The secondary system is the **domain name system (DNS)** which uses natural language **DNS names**.
- IP addresses are 32-bit integers.
  - ▶ Composed of four 8-bit octets.
  - ▶ Represented as four integers, usually in base 2 or base 10, separated by dots.

base 2: 11000111.00010001.00101000.11010010.  
base 10: 199.17.40.210.

# IP Address Assignment

- IP addresses are assigned to **network interfaces, not hosts**.
  - ▶ A host with one network interface is assigned an IP address by default.
- A network interface is normally assigned a unique IP address.
  - ▶ In practice, an interface may be assigned more than one address.
  - ▶ In some rare cases, an interface may be assigned no address at all.
  - ▶ In certain cases, different interfaces may have the same IP address.

# Class Networks (1)

- IP addresses are organized into **class networks** to facilitate address assignment and packet routing.
  - ▶ Class A: 0nnnnnnn.iiiiiii.iiiiiii.iiiiiii.
  - ▶ Class B: 10nnnnnnn.nnnnnnnn.iiiiiii.iiiiiii.
  - ▶ Class C: 110nnnnn.nnnnnnnn.nnnnnnnn.iiiiiii.
  - ▶ Class D: 1110bbbb.bbbbbbbb.bbbbbbbb.bbbbbbbb.
- A class network is really a set of IP addresses and not a network.
- The **Internet Corporation for Assigned Names and Numbers (ICANN)** is responsible for assigning class A, B, and C networks to organizations.
- Each address in a class A, B, or C network is a pair  $(N, I)$  where  $N$  is its **network identification** and  $I$  is its **interface (or host) identification**.  $n$  and  $i$  denote bits in the network and interface identifications, respectively.

## Class Networks (2)

- By convention, the address in a class A, B, or C network whose interface bits are all 0 (e.g., 199.17.40.0) is the **network address** for the class.
- By convention, the address in a class A, B, or C network whose interface bits are all 1 (e.g. 199.17.40.255) is the **(direct) broadcast address** for the class.
- The **limited broadcast address** is 255.255.255.255.
- The addresses in a class D network are for multicasting.
- Some addresses in the above classes and the addresses of the form 1111bbbb.bbbbbbbb.bbbbbbbb.bbbbbbbb are reserved.

# The Loopback

- Each host running TCP/IP has a virtual interface called the **loopback interface** which is the only interface on a virtual network called the **loopback network**.
- The network and interface addresses of the loopback are 127.0.0.0 and 127.0.0.1, respectively.

# Weaknesses of IP Address System

- Some hosts (e.g., a multi-homed host) have more than one IP address.
- The class networks are too rigid.
- There are not enough IP addresses for future expansion.

# Subnets

- IP addresses are also organized into **subnets** to facilitate address assignment, network organization, and routing.
- Each subnet is a set of addresses determined by:
  1. A **subnet address** (e.g., 199.17.35.96).
  2. A **subnet mask** (e.g., 255.255.255.240).
- Each address in a subnet is pair  $(S, I)$  where  $S$  is its **subnet identification** and  $I$  is its **interface (or host) identification**.
- Special cases:
  - ▶ Set of all IP addresses.
  - ▶ Class A, B, and C networks.
  - ▶ Individual interface (or host) IP address.



# Subnet Conventions

- Usually, but not necessarily, the subnet identification of a class A, B, or C address is an extension of the network identification of the address.
- Usually, but not necessarily, the subnet mask consists of a block of 1s followed by a block of 0s.
- By convention, there is one subnet corresponding to each SPN.
  - ▶ Each interface on the SPN is assigned the same subnet address and subnet mask.

# Address Resolution Problem

- High-level IP addresses are used for communication across an internet and are assigned independently of physical hardware addresses.
- Low-level physical addresses are needed for physically delivering a packet to an interface on a network.
- How are IP addresses mapped to physical addresses?
  - ▶ A solution is an **address resolution function**  $f$  that maps each IP address  $i$  to a physical address  $f(i)$ .
  - ▶ The function must be changed as the internet changes.
  - ▶ The function must be represented efficiently.

# Address Resolution Solutions

1. Physical addresses are encoded in IP addresses.
  - ▶ Possible for proNET networks.
  - ▶ Not viable for Ethernet.
2. Each machine contains a table that represents the local part of an address resolution function.
  - ▶ Awkward for Ethernet because physical addresses change when a host or interface is replaced.
3. IP addresses are bound to physical addresses dynamically.

# Address Resolution Protocol (ARP)

- Used for dynamically binding an IP address to a physical address (especially on Ethernet networks).
- ARP process:
  1. A host  $h_A$  broadcasts a request for the physical address which resolves an IP address  $i$ .
  2. The host  $h_B$  with the network interface having the address  $i$  sends a reply to  $h_A$  containing the physical address of the interface.
- The results of ARP queries are kept in a cache on each host.
- When a sender requests a physical address, it can include its physical address in the message.

# Complications

- Several packets may simultaneously need to know the same physical address.
- The host of the requested physical address may be down.
- The cache may contain out-of-date bindings.
  - ▶ When a host boots it can send a broadcast message informing the other computers on the network of its physical address.
- Hosts may provide bogus address bindings.
- At boot-time a diskless host knows its physical address but not its IP address.
  - ▶ The host must get its IP address from a server on another computer.

# Reverse Address Resolution Protocol (RARP)

- Used for obtaining the IP address that is bound to a physical address.
- RARP process:
  1. A host  $h_A$  broadcasts a request for the IP address which reversely resolves a physical address  $p$ .
  2. The RARP servers which receive the request send replies back to the  $h_A$  containing the requested IP address.
- Some scheme is needed to keep all the RARP servers from sending replies at the same time (and causing collisions on an Ethernet network).

# ARP and RARP Messages

- ARP and RARP messages are **encapsulated** in a physical frame.
  - ▶ ARP and RARP share the same message format.
  - ▶ Type field says the data is an ARP or RARP message.
  - ▶ The message itself is held in the data portion of the frame.
- Each message has the following address fields:
  - ▶ Sender IP address.
  - ▶ Sender physical address.
  - ▶ Target IP address.
  - ▶ Target physical address.