

SE 4C03 Winter 2007

11 Routing Protocols

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Routing Management Goals

- Packets can get from any source to any destination.
- Packets take the shortest route.
- Routing information is automatically distributed.
- Routing tables are automatically initialized and updated.
- Many routers route with partial information about destinations by using default routes.
 - ▶ Tables are smaller and routing is more efficient.
 - ▶ Management can be performed locally.
 - ▶ Routing tables are less likely to be inconsistent with each other.

Core Router Architecture

- The pool of internet routers is divided into **core routers** and **noncore routers**.
 - ▶ Motivated by the ARPANET which was a single WAN.
 - ▶ Works well when the internet has a single backbone.
- Provides optimal routes for all possible destinations.
 - ▶ Core routers do not use default routes.
 - ▶ Every assigned class network address must be advertised to the core system.
 - ▶ A noncore router sends all nonlocal traffic to the core router at its site.
 - ▶ Core routers communicate with themselves to preserve consistency.
- Core router architecture is impractical today because:
 - ▶ The Internet is no longer built around one backbone.
 - ▶ It is not possible to have a core router at each site.
 - ▶ Core router architecture does not scale up very well.

Peer Backbone Architecture

- Peer backbone networks have two or more backbone networks with several connections between them.
 - ▶ Motivated by connection of the NSFNET backbone to the ARPANET backbone.
- The system is difficult to implement because:
 - ▶ Packets cannot simply be routed according to the network portion of their destination.
 - ▶ Peers must keep their routes consistent with each other.
 - ▶ Default routes from one peer to another can create routing loops.

Vector Distance Routing (1)

- Each router keeps a list of route records having the following fields:
 - ▶ The destination class network (i.e., the **vector**).
 - ▶ The number of hops to the destination (i.e., the **distance**).
 - ▶ The route (either **direct** or the name of a **router**).
- **Assumption:** Measuring distance as number of hops is a good measure of the time cost of a route.
- When a router boots, the list is initialized to just the routes for the class networks that are directly connected to the router.
- Each router sends a copy of its list to all other routers that are directly connected to it.

Vector Distance Routing (2)

- Suppose that R_1 and R_2 are directly connected routers.
 - ▶ If R_1 has a route record (N, D, R) for the class network N but R_2 does not, R_2 will add the record $(N, D + 1, R_1)$ to its list.
 - ▶ If R_1 and R_2 have route records (N, D_1, R'_1) and (N, D_2, R'_2) , respectively, with $D_1 + 1 < D_2$, then R_2 will update its record to $(N, D_1 + 1, R_1)$.
 - ▶ If R_1 and R_2 have route records (N, D_1, R) and (N, D_2, R_1) , respectively, with $D_1 + 1 \neq D_2$, then R_2 will update its record to $(N, D_1 + 1, R_1)$.

Vector Distance Routing (3)

- **Advantage:** Vector-distance algorithms are easy to implement.
- **Disadvantage:** Vector-distance algorithms do not scale up well.
 - ▶ Each router eventually has a record in its list for each class network.
 - ▶ Lots of routing information is transmitted.
 - ▶ Routing information is propagated slowly (which can make the system unstable).
 - ▶ Any incorrect routing information will be propagated along with correct information.

Gateway-to-Gateway Protocol (GGP)

- Now defunct vector-distance protocol used for sharing routing information among core routers.
- GGP messages were encapsulated in IP datagrams (with protocol field set to 3).
- Kinds of GGP messages:
 - ▶ Routing update message.
 - ▶ Positive acknowledgment to routing update (update acceptable).
 - ▶ Negative acknowledgment to routing update (error detected).
 - ▶ Echo request.
 - ▶ Echo reply.

Link-State Routing (1)

(Shortest Path First (SPF) Routing)

- Each router keeps a graph of the topology of the internet.
 - ▶ A node represents a router.
 - ▶ An edge represents a network.
- The routers work to keep the graphs up to date.
 - ▶ Each router periodically checks to see if its neighbors are up or down.
 - ▶ Each router periodically broadcasts a message that contains the state of each of its links.
 - ▶ Each router uses the link-state broadcasts it receives to update its internet topology graph.

Link-State Routing (2)

- Whenever the graph changes, a router uses the **Dijkstra shortest path algorithm** to compute the shortest route to each destination.
- Advantages over vector-distance routing.
 - ▶ Each router independently computes the shortest routes from the same information.
 - ▶ It is easy to fix mistakes because link-state information is not modified as it is propagated.
 - ▶ The size of link-state messages do not grow as the internet grows.

Problems with Core Router System

- It is impractical for the group of core routers to include more than a small portion of the internet routers.
- Noncore routers need to know core router routes to avoid “extra hops” .
- Core routers need to know about “hidden” networks.

Autonomous System Architecture

- An internet is composed of several **autonomous systems** each composed of a collection of routers and networks.
- Each autonomous system chooses its own internal routing architecture.
 - ▶ An **interior gateway protocol (IGP)** is used to distribute routing information within an autonomous system.
- One or more routers in an autonomous system advertises local routing information to other autonomous systems.
 - ▶ An **exterior gateway protocol (EGP)** is used to distribute routing information between two autonomous systems.

Border Gateway Protocol (BGP)

- BGP is the EGP currently used in the Internet and most other TCP/IP internets.
 - ▶ Each autonomous system has at least one designated **border gateway** that speaks on behalf of the autonomous system.
 - ▶ Border gateways exchange reachability information (but not optimal routes).
- Attributes of BGP:
 - ▶ Transport is via TCP and thus is reliable.
 - ▶ Local routing policy can be supported.
 - ▶ Updates are usually incremental.
 - ▶ A receiver can authenticate the sender.
 - ▶ Subnet addressing is supported and so, for example, related destinations can be aggregated.

BGP Message Types

1. OPEN: initializes communication.
2. UPDATE: withdraws unreachable destinations and advertises new destinations. Each advertised destination can include:
 - ▶ The next hop to the destination.
 - ▶ The path of autonomous systems to the destination.
 - ▶ The source of the advertisement.

Each destination is expressed as a compressed address-mask pair.

3. NOTIFICATION: reports an error.
4. KEEPALIVE: tests network connectivity.
 - ▶ KEEPALIVE messages have minimum size (they contain only a BGP header).

Routing Information Protocol (RIP) (1)

- An IGP protocol for vector-distance routing.
 - ▶ Distance is measured by the number of hops.
- Implemented by the `routed` program designed at the University of California at Berkeley.
 - ▶ Became popular because it was distributed with BSD UNIX.
- RIP participants are either **active** or **passive**.
 - ▶ Active participants are routers that advertise their routes to others.
 - ▶ Passive participants are hosts that do not advertise routes but use advertisements to update their routes.

Routing Information Protocol (RIP) (2)

- An active participant broadcasts a message every 30 seconds.
 - ▶ RIP messages are encapsulated in UDP datagrams.
 - ▶ A RIP server listens at UDP port 520.
- Route restrictions.
 - ▶ An old route is retained until a new one with a strictly lower cost is received.
 - ▶ Routes timeout after 180 seconds.
 - ▶ The maximum possible distance is 16, so RIP only works with relatively small autonomous systems.
- RIP suffers from the problems that are inherent in vector-distance routing.
- RIP does not employ router authentication.

The HELLO Protocol

- A defunct IGP protocol for vector-distance routing.
- Routes are measured by network delay instead of number of hops.

The Open SPF Protocol (OSPF) (1)

- An IGP protocol for link-state (SPF) routing.
- Provides support for:
 - ▶ **Type of service routing**: Routing is done on the basis of both destination address and type of service (precedence plus low delay, high throughput, or high reliability).
 - ▶ **Load balancing**: The same traffic can be distributed over multiple routes.
 - ▶ **Area organization**: Local networks and routers can be organized into independent **areas**.
 - ▶ **Authentication**: Routers must authenticate each other.
 - ▶ **Subnet routes**: Routes may be directed to subnets.
 - ▶ **Information forwarding**: Routers may forward information received from routers exterior to a site.

The Open SPF Protocol (OSPF) (2)

- Each router keeps a graph of the topology of the internet whose structure is more complex than that of other link-state routing protocols.
 - ▶ A node represents a router or an SPN.
 - ▶ One router can be directly connected to another router in the graph.
- Kinds of OSPF messages:
 - ▶ Hello (to test router reachability).
 - ▶ Graph description (to initialize a router's topology graph).
 - ▶ Link status (to update a link in a router's topology graph).
- OSPF messages are encapsulated in IP datagrams (with protocol field set to 89).

The gated Program

- Handles multiple routing protocols including both IGPs and EGPs.
- Allows a router to communicate with routers both inside and outside its autonomous system.
- IGPs supported include: RIP, HELLO, OSPF.
- EGPs supported include: BGP.